

while $\text{sizes}(L) < 0 : L = -\text{sizes}(L)$

Write L to the LABEL lattice, $\text{sizes}(L) \neq +1$

- both are occupied:

find both root labels, take the smaller and unify them

$$\left(\begin{array}{l} L_1 < L_2 : \text{sizes}(L_1) \pm \text{sizes}(L_2) + 1 \\ \text{sizes}(L_2) = -L_1 \end{array} \right)$$

at the end the SIZES array has what we need.

~~fill label~~ TRICK: do this alg as occupations are generated. we can forget old rows. (we need L^{d-1} storage)

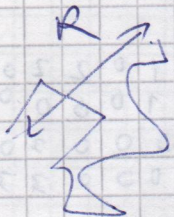
$$\left(\begin{array}{l} \text{in } d=2 \quad p = \frac{1}{2} \text{ for triangular latt} \\ \quad \quad \quad p = 0.5927 \text{ square latt} \end{array} \right) \quad \nu = 4/3 \quad \beta = \frac{5}{36}$$

for finite lattice e.g. $P_{\infty} = \frac{\text{points in largest cluster}}{L^d}$

(generations bond percolation directed percolation) P_{∞} L^d $L=200$ $L=100$ $L=50$

Random walks

Conformation of a polymer chain



$$R \sim N^{\nu}$$

Simplest model: Random walk

$\rightarrow R \sim t^{1/2}$ disagrees with experiment.

no back track RW: do not immediately turn back

SAW self avoiding random walk: do not go to a site already visited (monomers repulse)

What should be the weight of a random walk?

RW: all walks of length N are equally probable

SAW: ————— as long as they are SAWs

⇒ generate RVs, throw away if they intersect
 works for ~~to~~ small N , otherwise most RVs are
 thrown away (attrition problem)

Reptation algorithm:

1. Assume we have a walk.
2. delete one point from a random end
3. add one point to a random end with random direction, accept the new walk if it is SAW
 (otherwise use old as new config.)

Pivot algorithm Use lattice symmetries

$$v = \frac{3}{4} \quad d=2$$

$$v = 0.59... \quad d=3$$

$$v = 1/2 \quad d=4$$

Biased sampling:

Check how many steps are available
 (no backstroke)

$$W = \prod_{i=1}^N \frac{l_i}{(2d-1)} \leftarrow \begin{array}{l} \text{number of steps} \\ \text{at step } i \end{array} \text{ possible}$$

Generalization: potential between monomers.
 now different walks have different weights.

Importance Sampling:

Usual question: $\langle F \rangle = \frac{1}{Z} \int F(p, q) e^{-\beta H(p, q)} d^N p d^N q$

$e^{-\beta H(p, q)} \approx 0$ for most configurations

$$\langle F \rangle = \int F(E) \cdot \rho(E) \cdot e^{-\beta E} dE$$

$\rho(E) \cdot e^{-\beta E}$ is sharply peaked

uniform random configurations have too little weight ($\rho \approx 0$)
 or too large energy $\rho(E) e^{-\beta E} \approx 0$

if we sample with prob. $e^{-\beta H(p, q)}$

$$\langle F \rangle = \frac{1}{N} \sum F(p_i, q_i)$$

Metropolis's algorithm

Markov chain

i - computer time

$$\phi_i \rightarrow \phi_{i+1} \rightarrow \phi_{i+2} \dots$$

prob. to go from ϕ to $\phi' = T(\phi'|\phi)$

T does not depend on i , only on ϕ not older $\phi_s \Rightarrow$

\Rightarrow Markov process

$P(Q)$ = prob. of state Q

$$\sum_{\phi'} T(\phi'|\phi) = 1$$

Master eq:
$$\frac{\partial P(Q,t)}{\partial t} = \sum_{Q'} W(Q \leftarrow Q') P(Q') - \sum_{Q'} W(Q' \leftarrow Q) P(Q)$$

discretised

$$P_j^{(n+1)} = T_{jk} P_k^{(n)}$$

$$P_i^{(n+1)} = \sum_{\phi'} T(\phi_i|\phi') P_i^{(n)}$$

$$P_j^{(n+1)} - P_j^{(n)} = T_{jk} P_k^{(n)} - P_j^{(n)}$$

In equilibrium
$$T(\phi) = \sum_k T_{jk} P_k^{(n)} - \sum_k T_{kj} P_j^{(n)}$$

Balance equation
no source or sink

$$T(\phi) = \sum_{\phi'} T(\phi|\phi') \pi(\phi') = \sum_{\phi'} T(\phi'|\phi) \pi(\phi)$$

In Matrix notation

$$T_{ik} \phi_k = \phi_i$$

eigen vector with e.v. = 1

$\forall i, \phi_j \in N$ $(T_{ij} > 0$ irred
 $\forall i, \phi_j \in N$ $\exists t_{ij} > 0$ aperiodic
 $\forall i \in N$ $(T^m)_{ii} > 0$

Let's choose $T(\phi|\phi)$ such that we have

detailed balance

$$T(\phi'|\phi) \pi(\phi) = T(\phi|\phi') \pi(\phi')$$

current is the same in both directions

$$T_{ij} \phi_j = T_{ji} \phi_i$$

consider the evolution of $P_j^{(1)} \rightarrow P_j^{(2)} \dots P_j^{(n)} \rightarrow P_j^{(n+1)}$

does it converge to π ?

$$\sigma_N = \sum_i |P_i^{(N)} - \pi_i|$$

$$\sigma_{N+1} = \sum_i |T_{ij} P_j^{(N)} - \pi_i| = \sum_i |T_{ij} P_j^{(N)} - T_{ji} \pi_j| =$$

$$= \sum_i \left| \sum_j (T_{ij} P_j^{(N)} - T_{ij} \pi_j) \right| \leq \sum_{ij} T_{ij} |P_j^{(N)} - \pi_j| = \sigma_N$$

$\left. \begin{array}{l} \sum_i P_i^{(N)} - \pi_i = 0 \\ T_{ij} > 0 \end{array} \right\} \Rightarrow$ equality only if $\sigma_N = 0$
 \Rightarrow convergence to unique eq. P

Metropolis's alg.

starting from state ϕ_0 . After $n-1$ steps, do

1. proposal $T_0(\phi' | \phi_{n-1})$
2. Acceptance probability $T_A(\phi' | \phi_{n-1})$
if accepted $\phi_n = \phi'$
else $\phi_n = \phi_{n-1}$

To satisfy detailed balance, we must have

$$T_0(\phi | \phi') \cdot T_A(\phi | \phi') \cdot \pi(\phi') = T_0(\phi' | \phi) T_A(\phi' | \phi) \cdot \pi(\phi)$$

$$\frac{T_0(\phi | \phi') T_A(\phi | \phi')}{T_0(\phi' | \phi) T_A(\phi' | \phi)} = \frac{\pi(\phi)}{\pi(\phi')}$$

$$\frac{T_A(\phi | \phi')}{T_A(\phi' | \phi)} = \frac{\pi(\phi)}{\pi(\phi')} \cdot \frac{T_0(\phi' | \phi)}{T_0(\phi | \phi')} = F(\phi | \phi')$$

we can choose e.g. $T_A = \begin{cases} F(\phi | \phi') & \text{if } F(\phi | \phi') < 1 \\ 1 & \text{otherwise} \end{cases}$

$$\text{or } T_A(\phi | \phi') = \frac{F}{1 + F}$$

usually $T_0(\phi | \phi') = T_0(\phi' | \phi)$ reversibility of proposal

$$\pi(\phi) = e^{-S[\phi]}$$

$$\frac{\pi(\phi)}{\pi(\phi')} = e^{-(S(\phi) - S(\phi'))} = e^{-\Delta S}$$

in this case acceptance prob = $\begin{cases} e^{-\Delta S} & \text{if } \Delta S > 0 \\ 1 & \text{if } \Delta S < 0 \end{cases}$

$T_0(\phi | \phi') = ?$ should be optimized

change too big \Rightarrow acceptance small
change too small \Rightarrow conf. space exploration slow

e.g. for Ising model

local change $s_i \rightarrow -s_i$ OK for most cases
cluster alg: change many spins at once except near critical phase transition